Lecture 17

Gale-Shapley (Deferred Acceptance) Algorithm

Announcements

- HW8 due on Wednesday.
- This week's lectures (including this one) are NOT on the final exam.

- EthiCS pre-recorded lectures (5 short videos) are fair game for the final exam.
- Final exam: Two pages front-and-back of handwritten notes allowed.

Recap: One way to greedy algorithms

Greedy algorithms

- Make a series of choices.
 - Choose this activity, then that one, ...
 - Never backtrack.
- Show (or hope) that your choice never rules out success.
 - At every step, there exists an optimal solution consistent with the choices we've made so far.
- At the end of the day:
 - you've built only one solution,
 - never having ruled out success,
 - so your solution must be correct.

Recap: A different approach to greedy

- Greedy algorithms
 - Make a series of choices.
 - Choose this activity, then that one, ...
 - Never packtrack.
 - Instead: At each step, free to revert any of the choices we've already made – as long as the solution is improving!

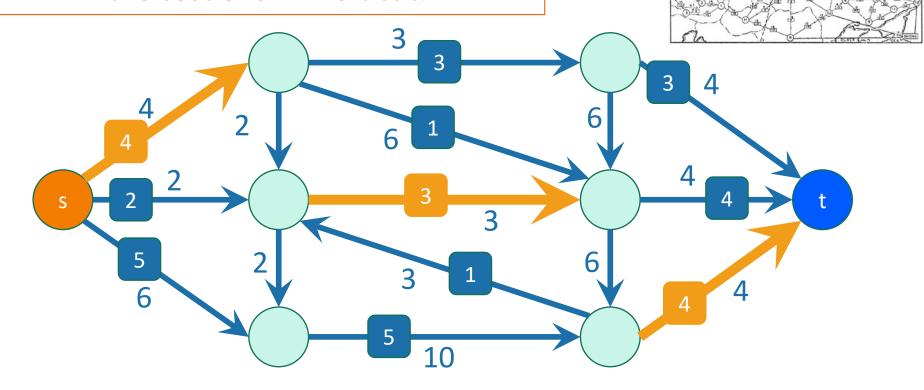
Recap: Ford-Fulkerson algorithm for s-t min-cut / max-flow

USA: s-t Min-Cut

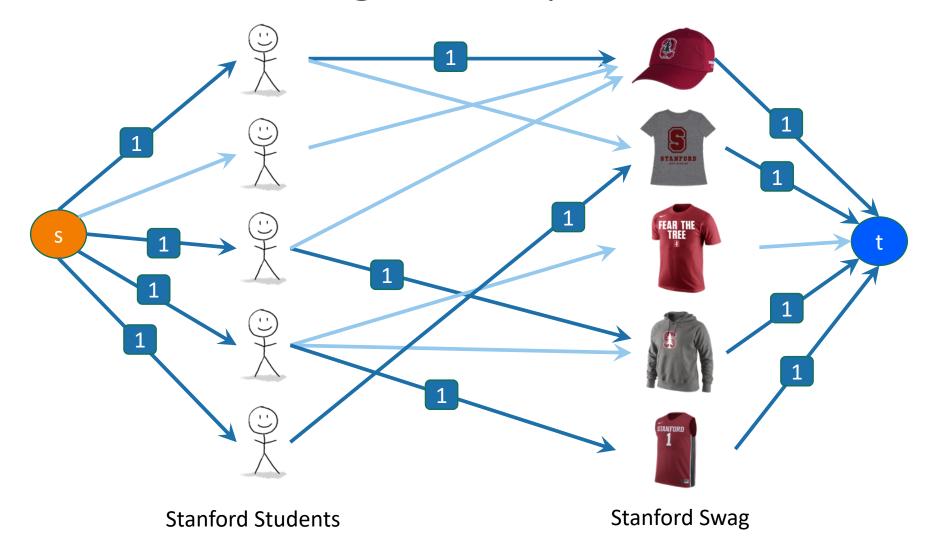
USSR: s-t Max-Flow

The value of a max flow from s to t

is equal to
the cost of a min s-t cut.



Recap: used s-t max-flow to solve assignment problems



Today: matching when both sides have preferences I want a CS161 CA to wear me!

Stanford Students

Stanford Swag

Today

• Hospitals/residents problem



- Stable matchings
 - Solve the hospitals/residents problem
 - But can we find them?

- Deferred Acceptance Algorithm
 - Find stable matchings!

Discussion, applications and non-applications

The hospital residency problem

- After completing medical school, students are finally ready to start their "residency" (similar to job internship):
 - In contrast, I'm told that many of you can get an internship after completing CS161...
- Each applicant has a preference over different residency programs.
- Each program has a preference over the applicants.
 How should you match applicants to residencies?

Simplifying assumption today: Each program has 1 slot

The hospital residency problem

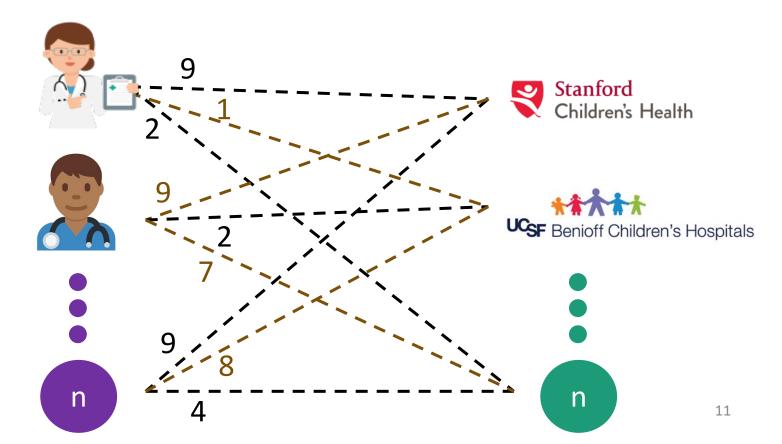
- After completing medical school, students are finally ready to start their "residency" (similar to job internship):
 - In contrast, I'm told that many of you can get an internship after completing CS161...
- Each doctor has a preference over different hospitals.
- Each hospital has a preference over the doctors.
 How should you match doctors with hospitals?

Simplifying assumption today: Each hospital has 1 slot

One way to model this problem

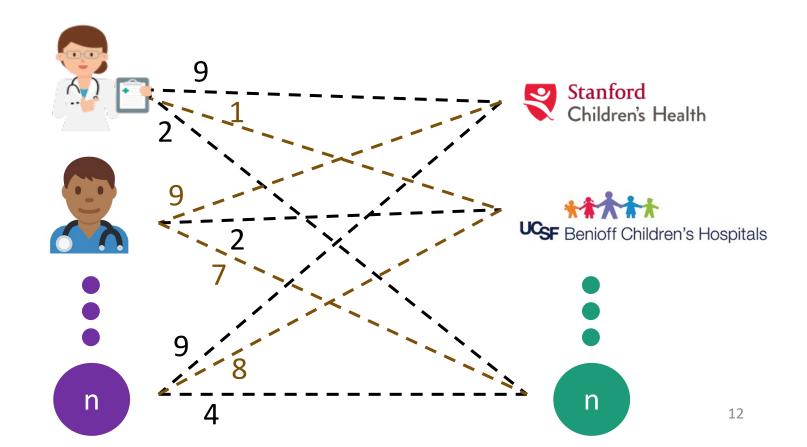
- Each doctor has a preference over hospitals
- Each hospital has a preference over the doctor

How should you match doctors with hospitals?



One way to model this problem

- Bipartite graph between doctors and hospitals
- Weights on edges = some function of preferences (highest weight = most preferred)

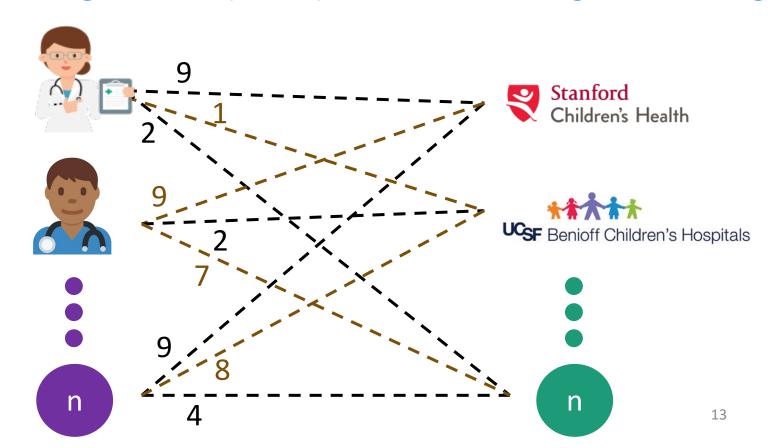


This slide just for intuition: You don't need to know Hungarian Algorithm!

One way to model this problem

- Bipartite graph between doctors and hospitals
- Weights on edges = some function of preferences

"Hungarian Algorithm" (CS261) finds a max weight matching

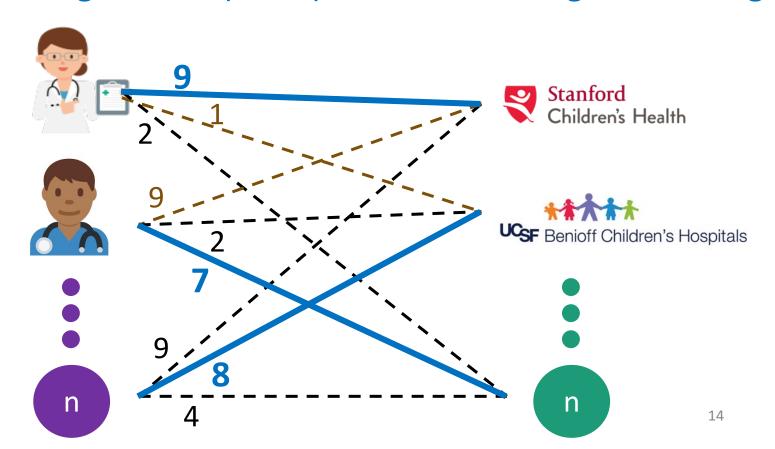


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One way to model this problem

- Bipartite graph between doctors and hospitals
- Weights on edges = some function of preferences

"Hungarian Algorithm" (CS261) finds a max weight matching



"Each hospital/doctor has a list of preferences"

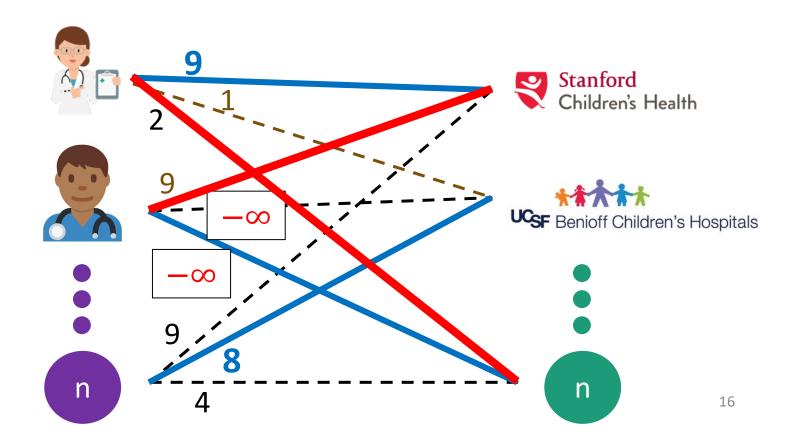
Missing step:

How does the *algorithm* get the preferences?

Where does your input come from?

... and what can go wrong if we don't think about it carefully:

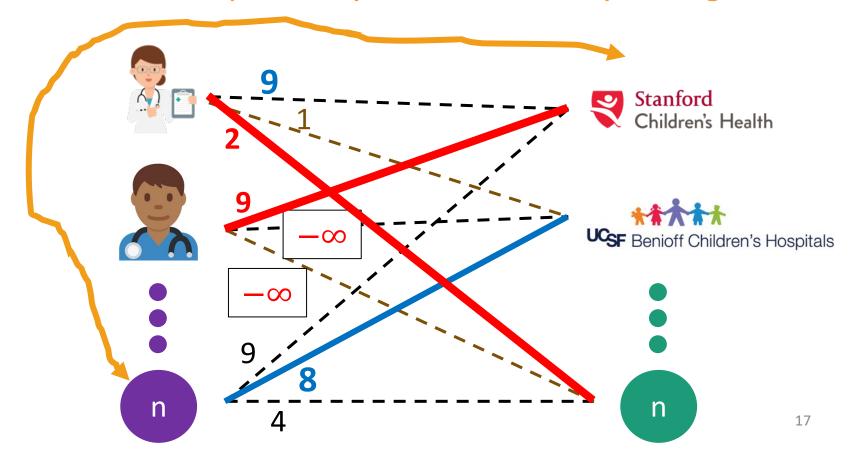
1. Some doctors may misreport their preferences



Where does your input come from?

... and what can go wrong if we don't think about it carefully:

- 1. Some doctors may misreport their preferences
- 2. Some doc+hospital may match outside your algorithm



Today

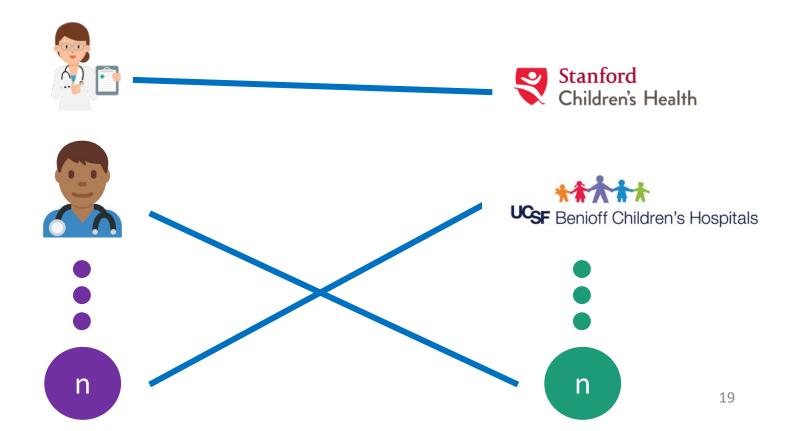
Hospitals/residents problem

Stable matchings

- Solve the hospitals/residents problem
- But can we find them?

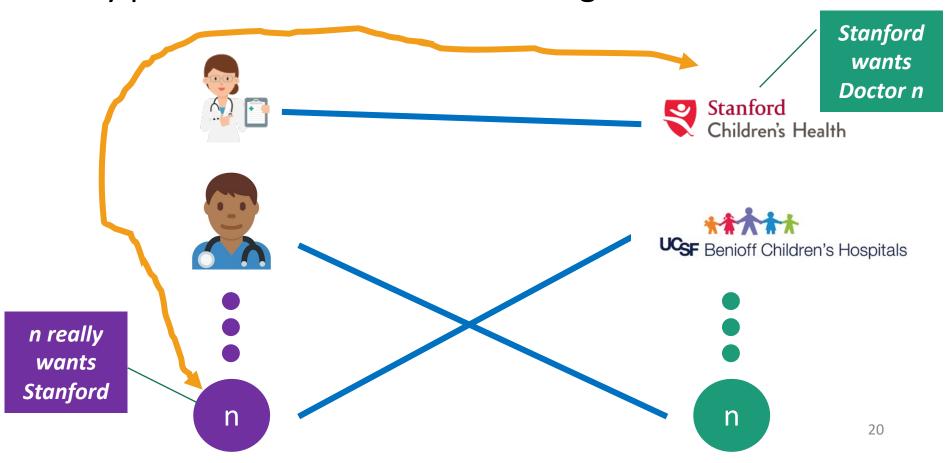
- Deferred Acceptance Algorithm
 - Find stable matchings!

Discussion, applications and non-applications



Definition (blocking pair):

Given Matching M, (Doctor i, Hospital j) are a blocking pair if they prefer each other to their assignment in M

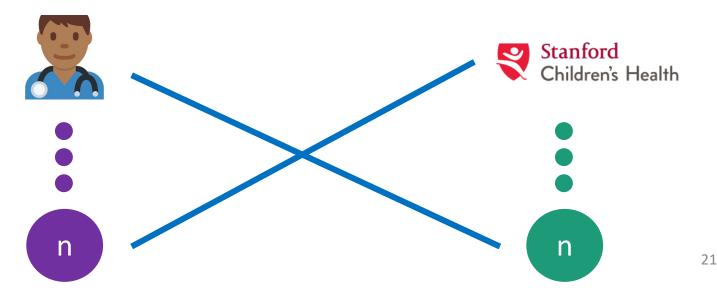


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Definition (stable matching):

M is a *stable matching* if there are no blocking pairs.



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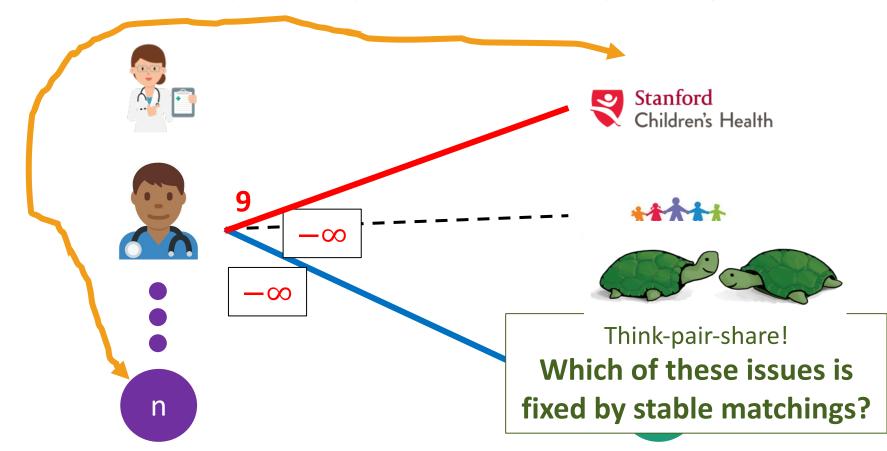
For every unmatched pair (i,j):

- Doctor i prefers Hospital M(i) over Hospital j, or;
- Hospital j prefers Doctor M(j) over Doctor i

Unstable Matching and incentives

Problems we identified with unstable matchings:

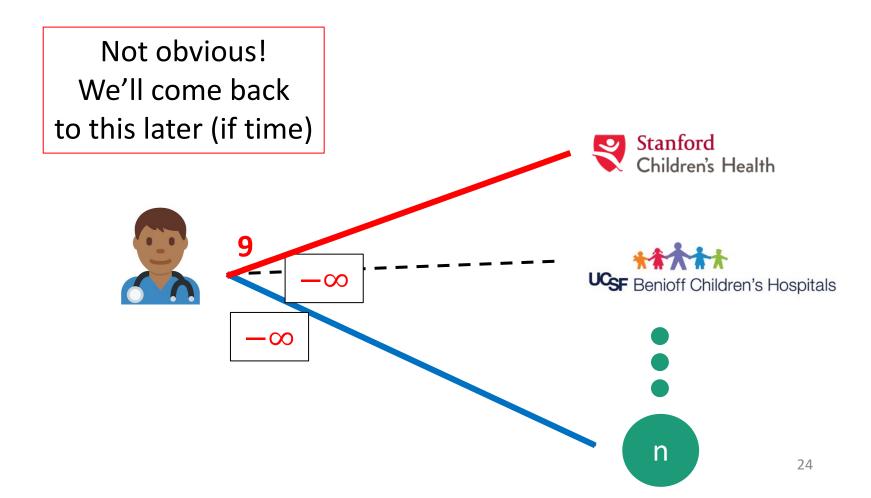
- 1. Some doctors may misreport their preferences
- 2. Some doc+hospital may match outside your algorithm



Stable Matching and incentives

With stable matching:

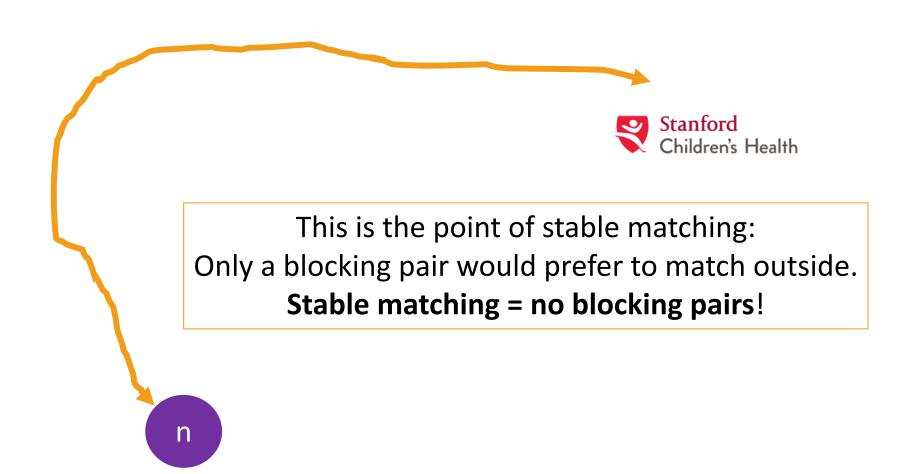
1. Will doctors misreport their preferences?



Stable Matching and incentives

With stable matching:

Doctor+hospital never prefer to match outside algorithm!



Stable Matching Problem

How to find stable matchings! (do they even exist?)

Stable Matching Problem

Stable Matching Problem

Input: each doctor/hospital submits a ranking (permutation) of {1,...,n}

Output: a stable matching

Alice's preferences		
1 st	Stanford	
2 nd	n	
•••		
n th	UCSF	

Stanford's preferences		
1 st	Alice	
2 nd	n	
	•••	
n th	Bob	

Definition (blocking pair):

Given Matching M, (Doctor i, Hospital j) are a *blocking pair* if they prefer each other to their assignment in M

Definition (stable matching):

M is a *stable matching* if there are no blocking pairs.

Naïve attempt #1

Greedy algorithm:

Step 1- match all the pairs (i,j) such that j is i's top choice, and i is j's top choice

Step 2- hopefully recurse on the rest somehow...

• Observation: Step 1 never rules out any solution ©

A slightly more ambitious attempt

Greedy attempt #2:

Step 1- try to match every doctor to her favorite hospital

Break ties by hospital preference

Step 2- hopefully recurse on the rest somehow...



A slightly more an

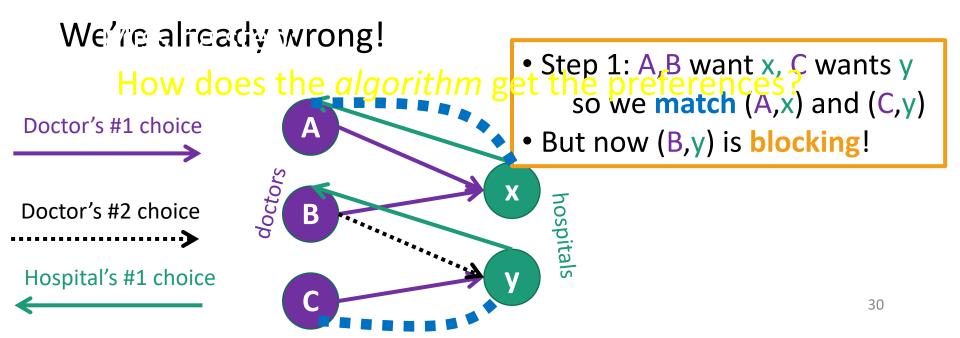
Think-pair-share!

Matching (C,y) was a bad idea...
How could we avoid it?

Greedy attempt #2:

Step 1- try to match every doctor to her favorite hospital

Break ties by hospital preference



Today

Hospitals/residents problem

- Stable matchings
 - Solve the hospitals/residents problem
 - But can we find them?

Deferred Acceptance Algorithm



• Find stable matchings!

Discussion, applications and non-applications

Questions?

Definition (blocking pair):

Given Matching M, (Doctor i, Hospital j) are a blocking pair if they prefer each other to their assignment in M

Definition (stable matching):

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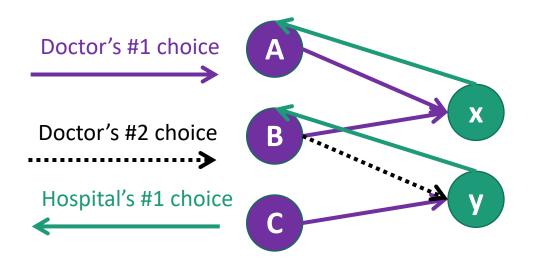
For every unmatched pair (i,j):

- Doctor i prefers Hospital M(i) over Hospital j, or;
- Hospital j prefers Doctor M(j) over Doctor i

[Gale Shapley '62] -> 2012 Nobel Prize* in Econ!

*- Joint w/ Al Roth from Stanford

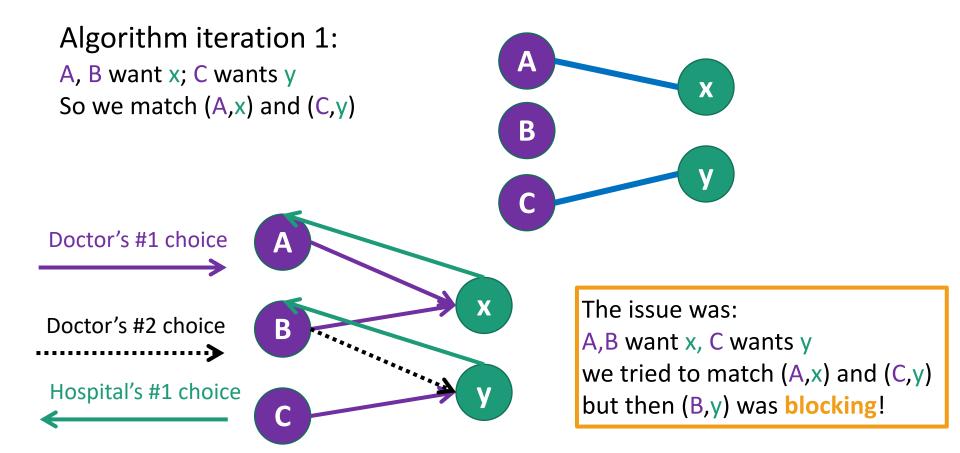
Main idea: *try* to match each doctor to top choice; if you discover a blocking pair, just switch the matching!



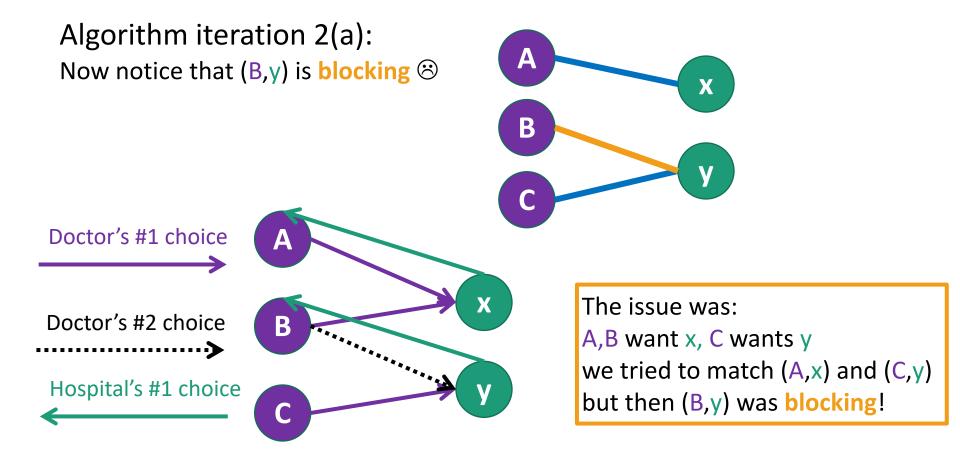
The issue was:

A,B want x, C wants y we tried to match (A,x) and (C,y) but then (B,y) was blocking!

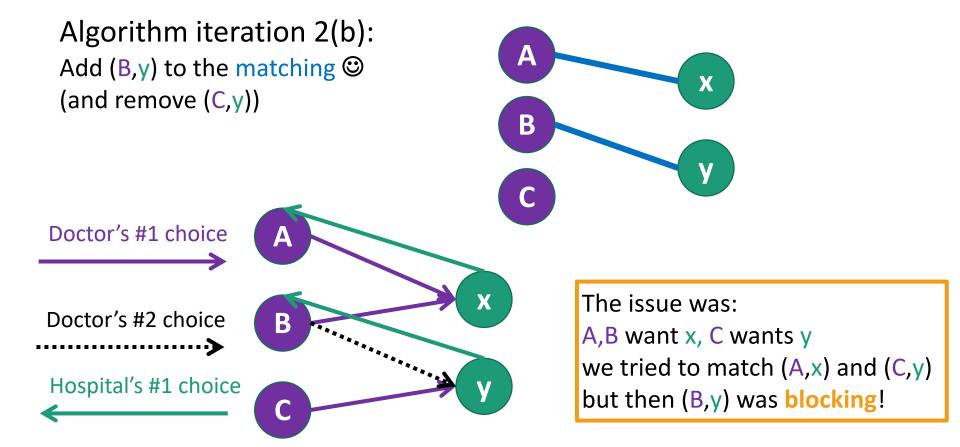
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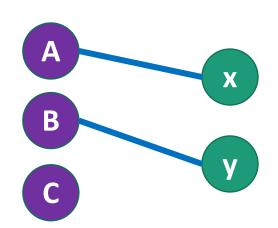
Main idea: *try* to match each doctor to top choice; if you discover a blocking pair, just switch the matching!

Algorithm iteration 2(b):

Add (B,y) to the matching (and remove (C,y))



Don't worry
Just switch around
until no blocking pairs!



The issue was:

A,B want x, C wants y we tried to match (A,x) and (C,y) but then (B,y) was blocking!

Lucky the Lackadaisical Lemur

Main idea: *try* to match each doctor to top choice; if you discover a blocking pair, just switch the matching!

Almost-pseudo-code:

While there is an unmatched doctor i:

Try to match i to next-favorite hospital on her list;

If this hospital doesn't have a doctor yet:

Both Doctor i and hospital are happy with this new match ©

Else-if this hospital prefers its current match i' over i:

Doctor i remains unmatched

Else-if this hospital prefers i over i':

Unmatch i'; Match (i, hospital)

Example run-through



X, Y, Z

Alice



Y, X, Z



Y, Z, X

B, A, C



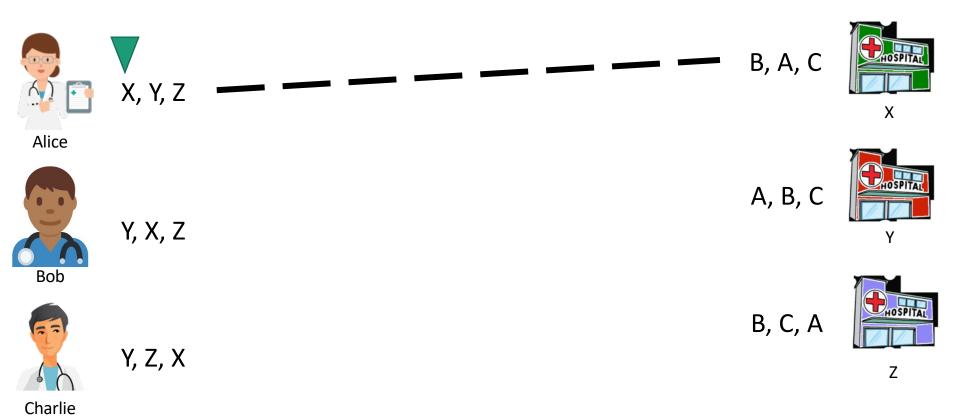
A, B, C

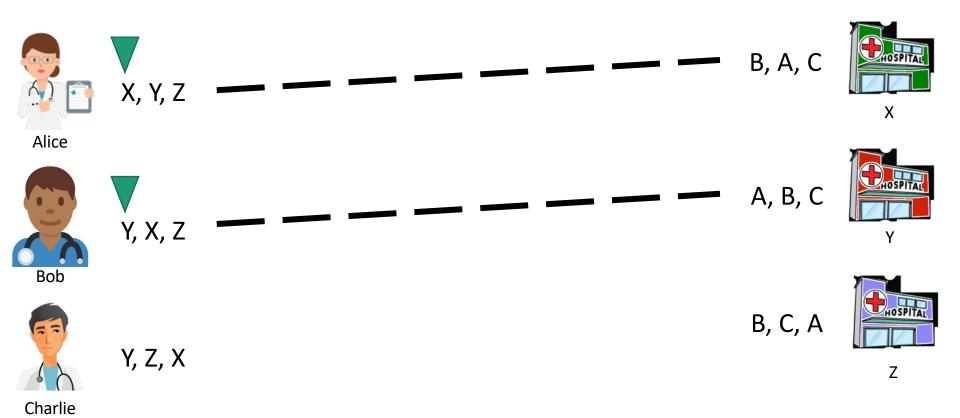


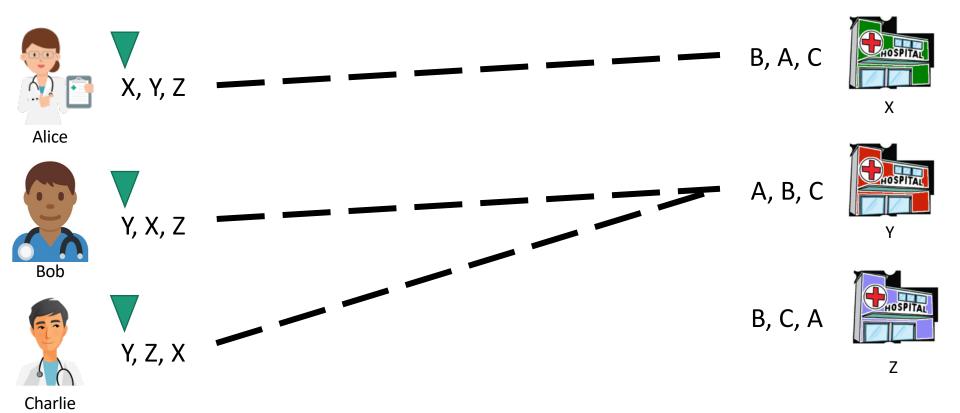
В, С, А

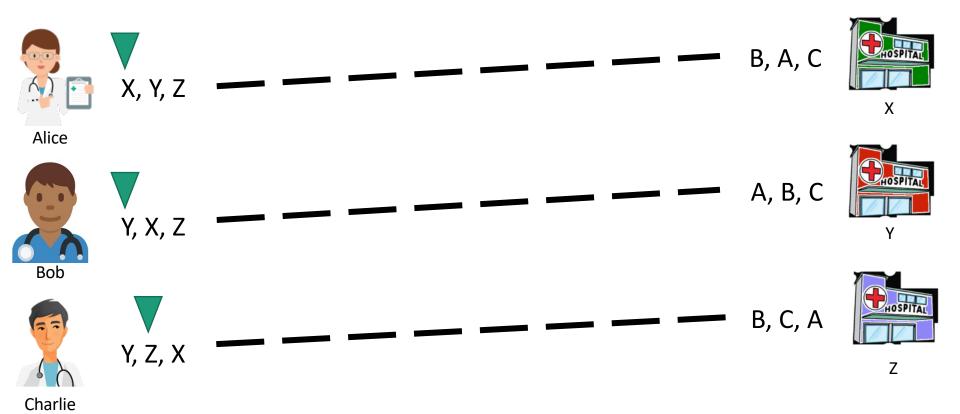


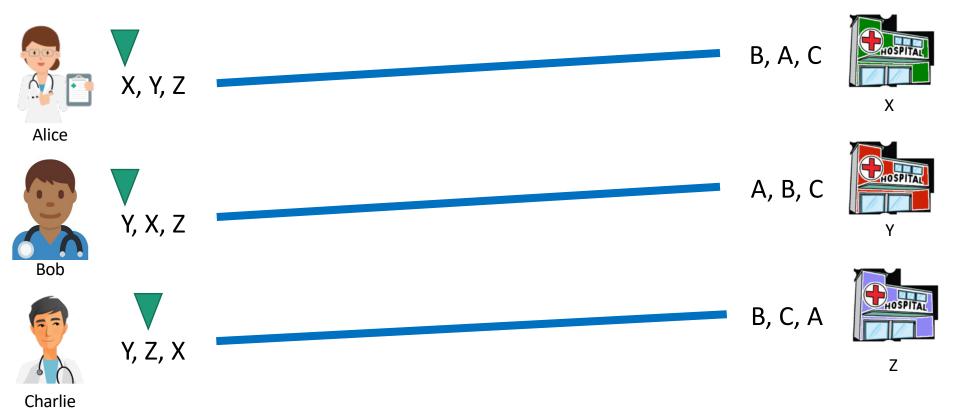
Ζ











Another example



X, Y, Z

Alice



Y, X, Z



Y, Z, X

В, А, С

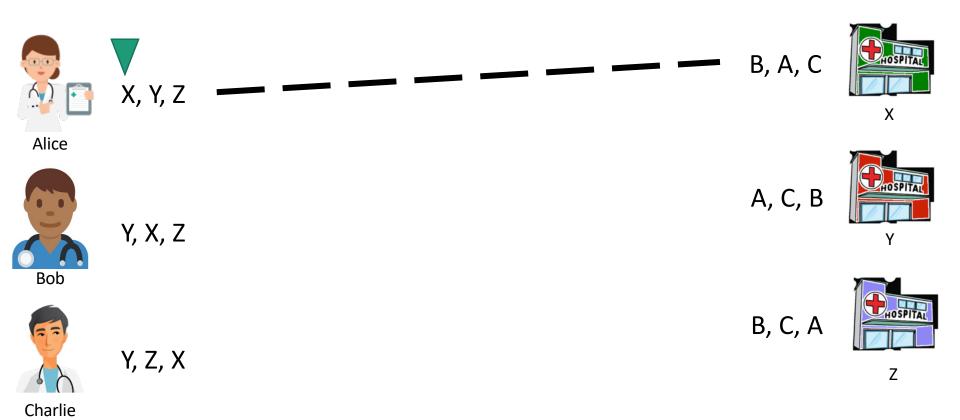


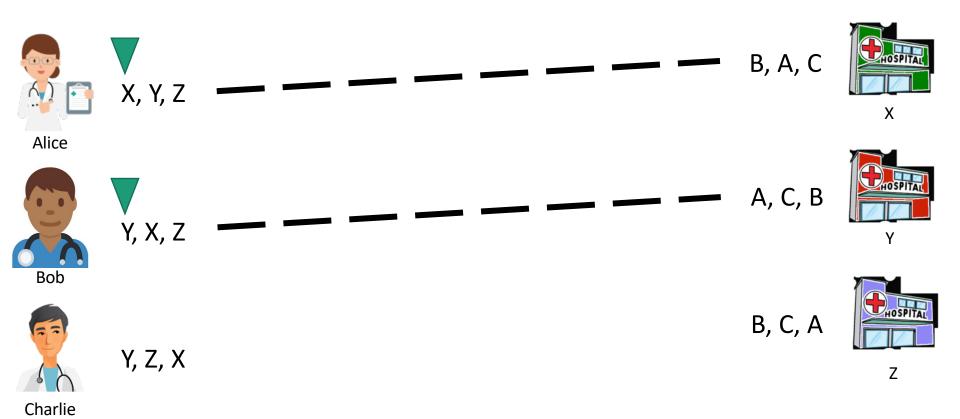
A, C, B

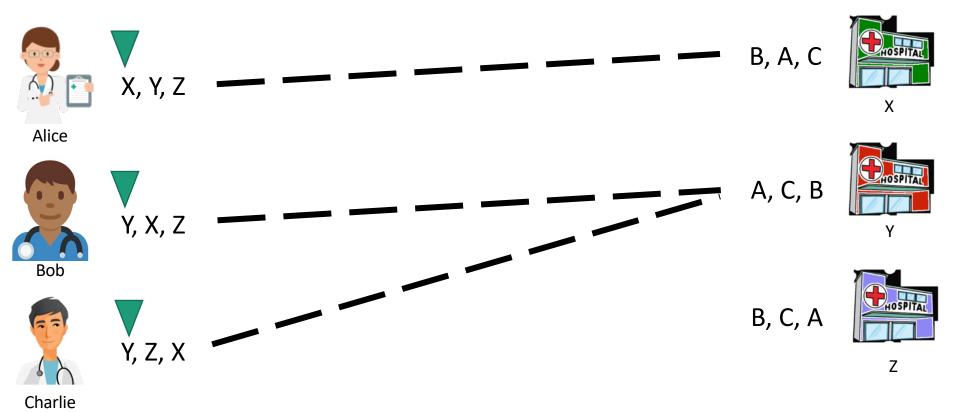


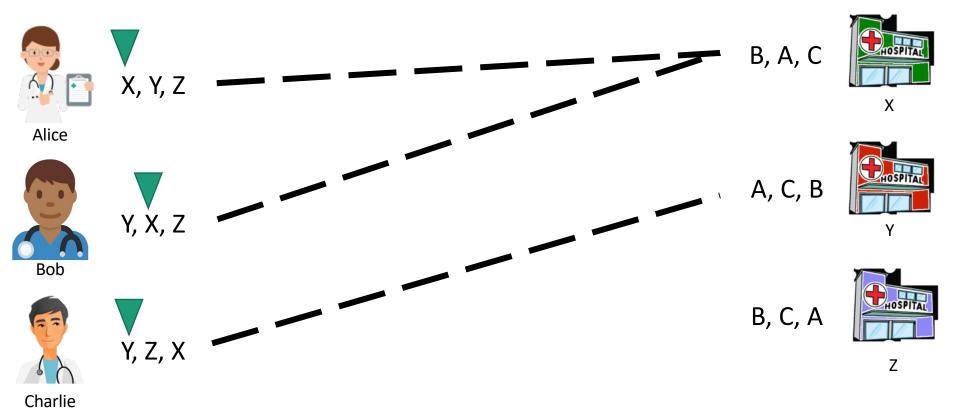


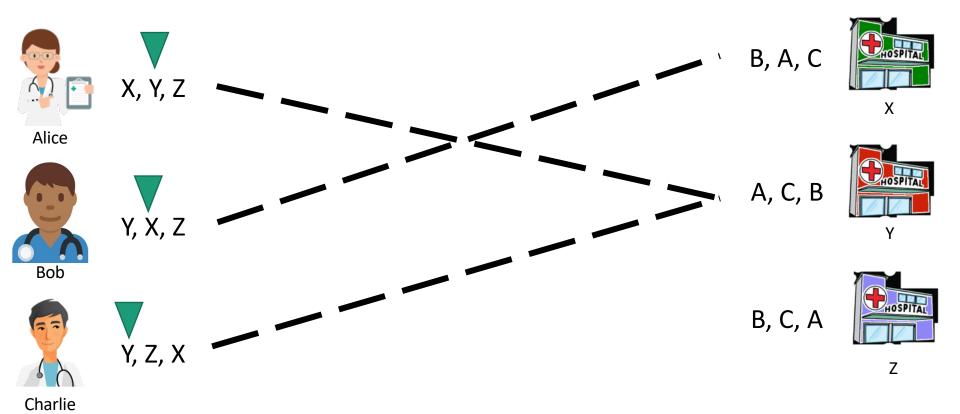
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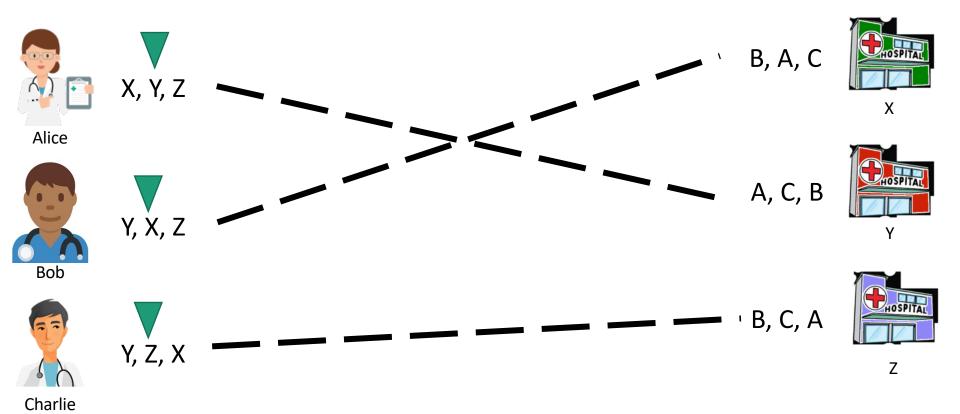


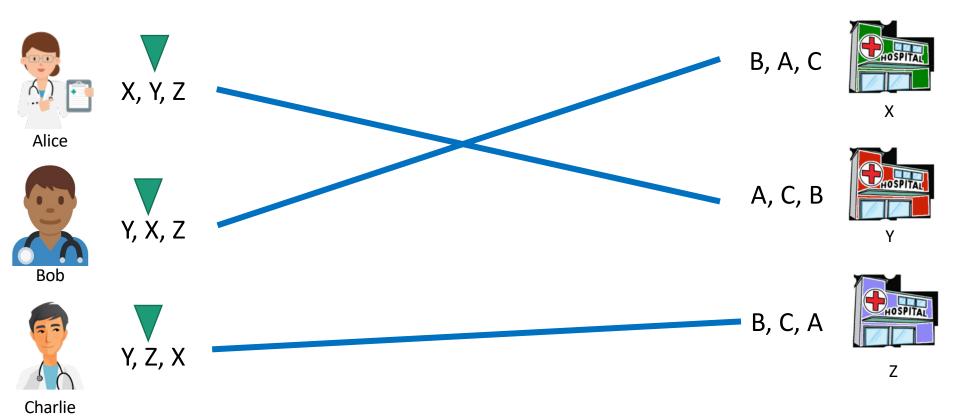












Deferred-Acceptance(Doctors, Hospitals):

```
// initialize

freeDoctors ← Doctors

for all d in Doctors:
    d.current ← 0

for all h in Hospitals:
    h.doctor ← NIL
```

```
// main loop
          while (exists d in freeDoctors)
               h ← d.ranking[d.current++]
                                              next favorite
               if (h is free)
                   h.doctor \leftarrow d
                   remove d from freeDoctors
               else-if (h.rank[d] < h.rank[h.doctor])</pre>
                   add h.doctor to freeDoctors
// h prefers d to
                   h.doctor \leftarrow d
previous match
                   remove d from freeDoctors
```

Think-pair-share! Running time?

return (h,h.doctor) for all h in Hospitals

Running time:

Each iteration of while loop = O(1)

Each iteration:

We +1 d.current for some doctor

We always have:

d.current $\leq n$ for every doctor (There are n doctors...)

Therefore, total run-time = $O(n^2)$

```
// main loop
while (exists d in freeDoctors)
    h ← d.ranking[d.current++]
                                    next favorite
    if (h is free)
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```

return (h,h.doctor) for all h in Hospitals

DA algorithm

Does it work?



• Yes!

- Is it fast?
 - O(n²) this is linear in the input size!

At worst exhaust through every doctor's preference list

Deferred Acceptance works!

Theorem: Given n doctors and n hospitals,

DA algorithm outputs a complete stable matching.

Corollary: A stable matching exists.

(This is not obvious!)

Proof of Theorem

Theorem: Given n doctors and n hospitals,

DA algorithm outputs a complete stable matching.

Proof: Follows from Claims 1+3 below...

<u>Claim 1:</u> At every iteration, current match is stable w.r.t. non-free doctors and hospitals.

<u>Claim 2:</u> Once a hospital is matched, it remains matched (possibly to a different doctor) until end of algorithm.

<u>Claim 3:</u> At the end of algorithm, every doctor/hospital is matched.

Proof of claims



<u>Claim 1:</u> At every iteration, current match is stable w.r.t. non-free doctors and hospitals.

Proof by contradiction: Suppose (d,h) blocking pair.

- \rightarrow d is currently matched to worse hospital than h.
- \rightarrow d already tried to match to h.
- \rightarrow h either refused d or left d later. Why?
- \rightarrow h must be matched to better doctor than d contradiction!

<u>Claim 2:</u> Once a hospital is matched, it remains matched (possibly to a different doctor) until end of algorithm.

"Proof": obvious from algorithm

<u>Claim 3:</u> At the end of algorithm, every doctor/hospital is matched.

Proof by contradiction: Suppose (d,h) still free.

End of algorithm $\rightarrow d$ already tried to match to h.

 \rightarrow after that step, h wasn't free \rightarrow by Claim 2, contradiction!

Deferred Acceptance works!

Theorem: Given n doctors and n hospitals,

DA algorithm outputs a complete stable matching.

Corollary: A stable matching exists.

<u>Claim 1:</u> At every iteration, current match is stable w.r.t. non-free doctors and hospitals.

<u>Claim 2:</u> Once a hospital is matched, it remains matched (possibly to a different doctor) until end of algorithm.

<u>Claim 3:</u> At the end of algorithm, every doctor/hospital is matched.

What have we learned?

Blocking Pair: A doctor and hospital that prefer each other over their respective matches.

Stable Matching: A matching without blocking pairs!

Deferred Acceptance Algorithm

"Tentatively match each free doctor to best interested hospital. Allow the hospital to leave match when a better doctor arrives."

Runs in time $O(n^2)$ = linear in input size \odot

Today

Hospitals/residents problem

- Stable matchings
 - Solve the hospitals/residents problem
 - But can we find them?

- Deferred Acceptance Algorithm
 - Find stable matchings!
- Discussion, applications and non-applications

The optimal stable matching?

DA algorithm found *a* stable matching...

- Is it *optimal*?
- What does optimality mean?



<u>Theorem:</u> The matching returned by DA is **doctor-optimal**, i.e. every doctor is matched to favorite hospital possible in any stable matching.

Corollary: Order of popping from freeDoctors does not change the output.

Theorem: Doctors cannot gain from misreporting their preferences.

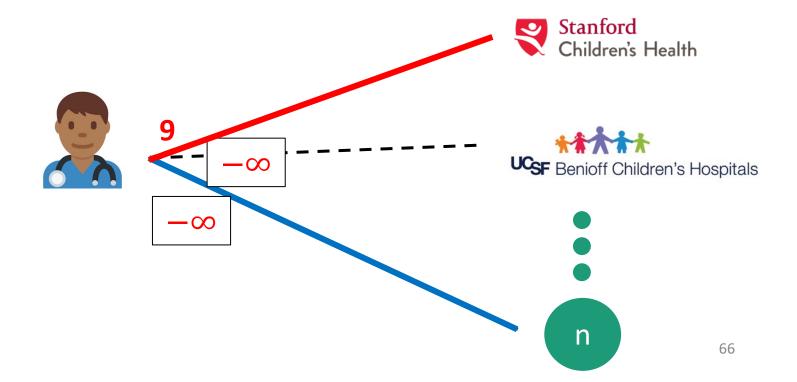


Stable Matching and Incentives

Doctor 2 may tell you he only wants to go to Stanford,

but...

Corollary: This won't help him if we find Stable Matching with DA!





The optimal stable matching?

<u>Theorem:</u> The matching returned by DA is **hospital-worst**, i.e. every hospital is matched to *least*-favorite doctor possible in any stable matching.

<u>Caution:</u> Hospitals *can* gain from misreporting their preferences.



Think-pair-share:

How would you find a hospital-optimal stable matching? Should actual matching be doctor- or hospital-optimal?

What have we learned?

<u>Doctor-optimality</u>: The matching returned by DA is <u>doctor-optimal</u> (but hospital-worst)

<u>Truthful preferences corollary:</u> Doctors cannot gain from misreporting their preferences (but hospitals *can*).

Point:

It's important to **think** about how **our algorithms affect people**. **Theorems** can help!

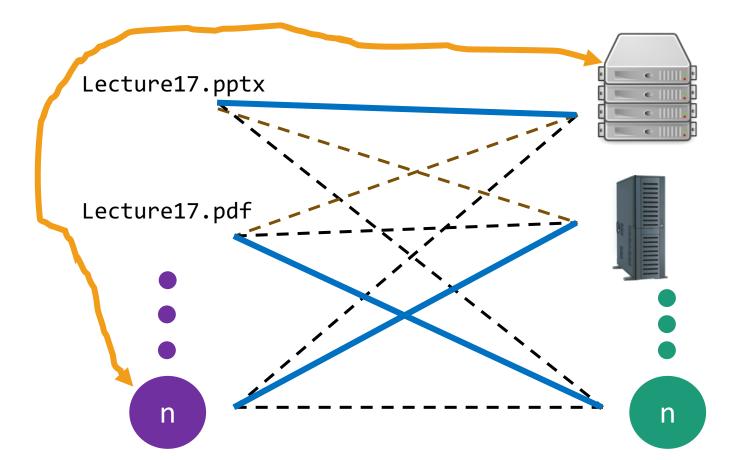
Today

Hospitals/residents problem

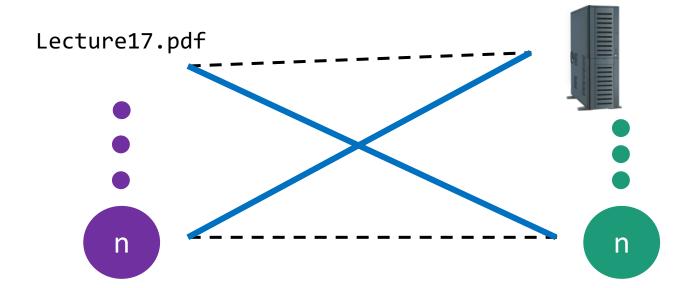
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• Suppose that instead of doctors and hospitals, you want to match packets to servers on the internet.



- Suppose that instead of doctors and hospitals, you want to match packets to servers on the internet.
- When you own all the servers, you don't have to worry about them matching outside your algorithm...
- But it turns out that Deferred Acceptance is just very fast in practice ©



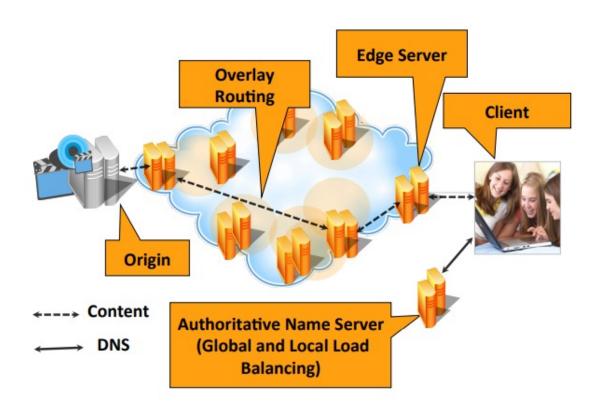
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Truncated preference lists

Packets typically get one of top servers

Total running time closer to O(n)!

Highly distributed: Every packet looks for its own server!



See "Algorithmic Nuggets in Content Delivery" (Maggs & Sitaraman, CCR'15) for details on how Akamai uses Deferred Acceptance to match packets to servers

Stanford Marriage Pact



Stanford Marriage Pact

- Matches between Stanford students who want to make a *pact*: "If we don't get married by time X, we'll marry each other."
- Historically, Gale-Shapley's original paper talked about Stable Marriage
 - men = doctors; women = hospitals.
- Original Marriage Pact used variant of Deferred Acceptance
 - It doesn't any more...



Marriage Pact doesn't need stability:
It is meant to be a back-up matchCouples are encouraged to find outside matches!

Recap

Hospitals/residents problem

- Stable matchings
 - Solve the hospitals/residents problem
 - But can we find them?

- Deferred Acceptance Algorithm
 - Find stable matchings!

Discussion, applications and non-applications

Next time

- Quick and hand-wavey recap of past lectures.
- Algorithms beyond 161 ...

