

# CS 161 W22: Section 3 Problems

## Exercise 0

### Warmup!

- True or False:* Expected runtime averages the runtime over the outcomes of random events within the algorithm and make no assumption about the input.
- I have an algorithm that takes positive integers  $(n, i)$  where  $1 \leq i \leq n$ . The algorithm rolls a  $n$ -sided die repeatedly until the die returns any value  $\leq i$ . What is the expected runtime in  $n$ ? Worst-case runtime? Rigorous proof not necessary :)

## Exercise 1

### Light Bulbs and Sockets

You are given a collection of  $n$  differently sized light bulbs that have to be fit into  $n$  sockets in a dark room. You are guaranteed that there is exactly one appropriately-sized socket for each light bulb and vice versa; however, there is no way to compare two bulbs together or two sockets together as you are in the dark and can barely see! You are, however, able to Spidey-sense where the sockets and light bulbs are, pick them up, and attempt to screw a light bulb into a chosen socket, from which you can determine whether the light bulb's base is too large, too small, or is an exact fit for the socket. If the bulb fits exactly, it will flash once, in which case you have a correct match. (Note that the flashing light does not allow you to visually compare bulbs/sockets to other bulbs/sockets.)

Suggest a (possibly randomized) algorithm to match each light bulb to its matching socket. Your algorithm should run strictly faster than quadratic time in expectation. Give an upper bound on the worst-case runtime, then prove your algorithm's correctness and expected runtime.

## Exercise 2

### All On the Same Line

Suppose you're given  $n$  distinct ordered pairs of integers  $(x_1, y_1), (x_2, y_2), \dots, (x_n, y_n)$ , where for all  $i, j$ ,  $x_i \neq x_j$  and  $y_i \neq y_j$ . Recall that two points uniquely define a line,  $y = mx + b$ , with slope  $m$  and intercept  $b$ . (Note that choosing  $m$  and  $b$  also uniquely defines a line.) We say that a set of points  $S$  is *collinear* if they all fall on the same line; that is, for all  $(x_i, y_i) \in S$ ,  $y_i = mx_i + b$  for fixed  $m$  and  $b$ . In this question, we want to find the maximum cardinality subset of the given points which are collinear – in less jargon, we're looking for the maximum integer  $N$  such that we can find  $N$  of the given points the same line. Assume that given two points, you can compute the corresponding  $m$  and  $b$  for the line passing through them in constant time, and you can compare two slopes or two intercepts in constant time.

*This is a challenging problem – so we're only going to pseudocode at a high level!*

- Design an algorithm to find a maximum cardinality set of collinear points in  $O(n^2 \log n)$  time. If there are several maximal sets, your algorithm can output any such set. *Some hints:*
  - $O(n^2 \log n) = O(n^2 \log n^2)$ , which looks like sorting  $n^2$  items.
  - Start small; how would we verify that 3 points are on the same line?
- It is not known whether we can solve the collinear points problem in better than  $O(n^2)$  time. But suppose we know that our maximum cardinality set of collinear points consists of exactly  $n/k$  points for some constant  $k$ . Design a randomized algorithm that reports the points in some maximum cardinality set in expected time  $O(n)$ . Prove the correctness and runtime of your algorithms.

*Some hints:*

- Your expected running time may also be expressed as  $O(k^2n)$ .
- Your algorithm might not terminate!

*For your own reflection:* Imagine that you, an algorithm designer, had to pick one of the algorithms in part (a) or (b) to implement in the autopilot of an airplane, as part of the route-planning of a self-driving car, or in any other scenario in which human lives are at stake. Given what you know about the performance and worst-case scenario of each of the algorithms, which algorithm would you choose and why?