1 Conditions for Shortest Path Algorithms

Suppose that we want to find the shortest path between two nodes in the following graph. Which algorithm can we use?



O BFS

O Dijkstra

O Bellman-Ford

O All of the above

O BFS and Dijkstra

O Dijkstra and Bellman-Ford

Correct

We have a graph with negative edge weights. Can we use Dijkstra to find shortest paths? O Yes

O No

Correct

We have an undirected graph with positive edge weights. Can we use Dijkstra to find shortest paths? Yes

O No

Correct

We have a directed graph with positive edge weights. Can we use Dijkstra to find shortest paths? Yes

O No

U NO

Correct

2 Dijkstra Forensics

Suppose we run Dijkstra on some graph with nodes A, B, C, D, E, F that has nonnegative (≥ 0) edge weights, starting from the node A. In the middle of the algorithm our computer crashes. We look through the memory dump, and see that the state of d looked as follows when the crash happened:

$$d[A] = 0, d[B] = 5, d[C] = 4, d[D] = 15, d[E] = 2, d[F] = 20$$

Additionally from the memory dump we see that the current node when the crash happened was node C.

What is the minimum possible length of the shortest path from node A to node B?



Correct

What is the maximum possible length of the shortest path from node A to node B?

5 Correct

What is the minimum possible length of the shortest path from node A to node D?

4

Correct

What is the maximum possible length of the shortest path from node A to node D?

Correct

15

What is the minimum possible length of the shortest path from node A to node E?



Correct

What is the maximum possible length of the shortest path from node A to node E?



What is the minimum possible length of the shortest path from node A to node F?



Correct

What is the maximum possible length of the shortest path from node A to node F?

Correct

If we run the Dijkstra algorithm on the graph of U.S. streets/roads/highways/etc., starting from the Stanford Oval, which of the following locations will become the current node first?

O Times Square in New York

O The Hollywood Sign

• Tresidder Union

O The ordering might differ in each run of Dijkstra.

Correct

3 Runtime

Suppose that we implement Dijkstra with a red-black tree. What is the asymptotically smallest upper bound on runtime in terms of n (the number of nodes) and m (the number of edges).

 $O O(n \log n + m)$

• $O((n+m)\log n)$ • O(n+m)

Correct

What if we implement Dijkstra with a Fibonacci heap? What is the asymptotically smallest upper bound on runtime in terms of n (the number of nodes) and m (the number of edges).

O(n log n + m)
O((n + m) log n)
O(n + m)

Correct

Suppose that we have a heap data structure that does not support updating the keys (many standard implementations of heaps in various programming languages do not support the update key operation).

Our data structure keeps a collection of items, each of form (key, object), where keys are numbers, and object can be anything (we will store vertices as our objects). Our data structure supports two operations:

- Insert a new (key, object) into the collection.
- Remove the item with the lowest key currently in the collection and return the key and object for it.

We run a modification of Dijkstra with the following pseudo-code:

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\begin{array}{l} d \leftarrow \text{array indexed with vertices and filled with } \infty \\ H \leftarrow \text{empty heap} \\ \text{Insert (0, starting node) into } H \\ \textbf{while } H \text{ is not empty } \textbf{do} \\ \hline \\ \textbf{Remove (key, vertex) from } H \text{ with the smallest key.} \\ \textbf{if } key < d[vertex] \textbf{ then} \\ \hline \\ d[vertex] \leftarrow key \\ \textbf{for all neighbors } w \text{ of vertex } \textbf{do} \\ \hline \\ \\ \textbf{Insert (} d[vertex] + weight(vertex, w), w) \text{ into } H. \end{array}
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What is the asymptotically smallest upper bound on the runtime of the above code assuming that both the insert and remove operations on *H* take $O(\log(\text{size of } H))$ time? Assume that $n - 1 \le m \le n^2$ (in particular log $m = \Theta(\log n)$).



- $O(m \log n)$
- O O(n+m)

Correct