CS 161
Design and Analysis of Algorithms

Lecture 1:
Logistics, introduction, and multiplication!

Slides originally created by Mary Wootters.
How was your break?
The big questions

• Who are we?
  • Professors, TAs, students?
• Why are we here?
  • Why learn about algorithms?
• What is going on?
  • What is this course about?
  • Logistics?
  • Embedded ethics?
• Can we multiply integers?
  • And can we do it quickly?
Who are we?

Instructors

Moses Charikar
Nima Anari

Awesome TAs

Amol Singh
Apoorva Dixit
Arpita Singhal
Jeff Z. HaoChen
Joey Rivkin

EthiCS Instructor

Course Manager

Dan Webber
Amelie Byun
Jonathan Ko
Junyao Zhao
Kamyar Salahi
Kayla Patterson
Mohammad Roghani

Paris Zhang
Rishu Garg
Samantha Liu
Thomas Mayer
Xinyi Wang
Who are you?

- Frosh
- Sophomores
- Juniors
- Seniors
- MA/MS Students
- PhD Students
- NDO Students

Concentrating in:

- Aero & Astro
- Archaeology
- Art History
- Asian American Studies
- Bioengineering
- Biology
- Biomedical Data Science
- Chemistry
- Civil & Env. Eng.
- Classics
- Computation and Mathematical Eng.
- Computer Science
- Creative Writing
- Data Science
- Earth System Science
- Economics
- Education
- Electrical Eng.
- Engineering
- English
- Epidemiology
- Human Biology
- Immunology
- Material Sci & Eng.
- Math
- Math & Comp. Sci.
- Mechanical Eng.
- Modern Languages
- MS&E
- Music
- Physics
- Political Science
- Psychology
- Sociology
- Symbolic Systems
- Theater & Performative Studies
- Undeclared
Where are you?
Why are we here?

• I’m here because I’m excited about algorithms!

Yay Algorithms!
Why are you here?

• Algorithms are fundamental.
• Algorithms are useful.
• Algorithms are fun!
• CS161 is a required course.

Why is CS161 required?

• Algorithms are fundamental.
• Algorithms are useful.
• Algorithms are fun!
Algorithms are fundamental

Operating Systems (CS 140)

Compilers (CS 143)

Networking (CS 144)

Cryptography (CS 255)

Computational Biology (CS 262)
Algorithms are useful

- All those things without the course numbers.
- As inputs get bigger and bigger, having good algorithms becomes more and more important!
Algorithms are fun!

• Algorithm design is both an art and a science.
• Many surprises!
• Many exciting research questions!
What’s going on?

• Course goals/overview
• Logistics
Course goals

• The design and analysis of algorithms
  • They go hand-in-hand

• In this course you will:
  • Learn to think analytically about algorithms
  • Flesh out an “algorithmic toolkit”
  • Learn to communicate clearly about algorithms
Roadmap

Today

- Divide and conquer

5 lectures
- Asymptotic Analysis
- Randomized Algs
- Recurrences

The Future!

- More detailed schedule on the website!

On FINAL
- On MIDTERM

8 lectures
- Longest, Shortest, Max and Min...
- Greedy Algs
- Dynamic Programming

2 lectures
- Data structures
- Randomized Algs
- Dynamic Programming
- Greedy Algs

2 lectures
- The Future!
Our guiding questions:

- Does it work?
- Is it fast?
- Can I do better?
Our internal monologue...

What exactly do we mean by better? And what about that corner case? Shouldn’t we be zero-indexing?

Does it work?
Is it fast?
Can I do better?

Dude, this is just like that other time. If you do the thing and the stuff like you did then, it’ll totally work real fast!

Plucky the Pedantic Penguin
Detail-oriented
Precise
Rigorous

Lucky the Lackadaisical Lemur
Big-picture
Intuitive
Hand-wavey

Both sides are necessary!
Aside: the bigger picture

• Does it work?
• Is it fast?
• Can I do better?

• Should it work?
• Should it be fast?

• We want to reduce crime.
• It would be more “efficient” to put cameras in everyone’s homes/cars/etc.

• We want advertisements to reach to the people to whom they are most relevant.
• It would be more “efficient” to make everyone’s private data public.

• We want to design algorithms, that work well, on average, in the population.
• It would be more “efficient” to focus on the majority population.
Course elements and resources

• Course website:
  • cs161.stanford.edu

• Lectures
  • Pre-Lecture Exercises
  • Lecture Notes
  • IPython Notebooks
  • Concept Check Questions

• References
• Sections
• Homework
• Exams
• Office Hours and Ed
Lectures

- Mon/Wed 10:30am-12:00pm (Skilling Auditorium)
  - Recorded on Panopto/Canvas
  - Fridays for EthiCS lectures and review sections
- Resources available:
  - Pre-lecture exercises
  - Slides, Videos, Notes, IPython notebooks, Concept Check Qns

Slides are the slides from lecture.

Videos from lecture are available!

Hand-outs and references have many details that slides may omit.

IPython notebooks have implementation details that slides may omit.
Embedded EthiCS Lectures

• Two lectures focused on ethics by Dan Webber: Jan 26 and Mar 1
  • Friday 10:30am-12:00pm (Skilling Auditorium)
• EthiCS questions will also be on your homework and your exams.
• More on EthiCS in a bit ...
How to get the most out of lectures

• **During lecture:**
  • Participate live (if you can), ask questions.
  • Engage with in-class questions.

• **Before lecture:**
  • Do *pre-lecture exercises* on the website.

• **After lecture:**
  • Go through the exercises on the slides.

• **Do the reading**
  • either before or after lecture, whatever works best for you.
  • *do not wait to “catch up” the week before the exam.*

Siggi the Studious Stork (recommended exercises)
Ollie the Over-achieving Ostrich (challenge questions)

Think-Pair-Share Terrapins (in-class questions)
IPython Notebooks

• Lectures will occasionally be accompanied by IPython notebooks (but not homework)
  • For the next lecture, the **pre-lecture exercise** is to get started with Jupyter Notebooks and with Python.
  • See the course website for details.

• The goal is to make the algorithms (and their runtimes) more tangible.
Concept Check Questions

• Not part of grade; will not be graded
• Links to question sets part of resources for each lecture (via Lectures tab on website)

Lecture resources

• Lecture notes: [PDF]
• Slides: [PDF] [PowerPoint]
• Python notebook: [Colab] [Zip]
• Concept check questions: [Interactive SVG] [Solved PDF]

Multiplication Algorithms

1 Grade-school multiplication

Suppose we multiply two \( n \)-digit integers \((x_1x_2\ldots x_n)\) and \((y_1y_2\ldots y_n)\) using the grade-school multiplication algorithm. How many pairs of digits \( x_i \) and \( y_j \) get multiplied in this algorithm?

- \( n^3 \)
- \( 2n - 1 \)
- \( n^2 \)
References

• **Algorithms Illuminated**, Vols 1, 2, and 3 by Tim Roughgarden

• Additional resources at algorithmsilluminated.org

• We may also refer to the following (optional) books:


  “Algorithm Design” by Kleinberg and Tardos
Sections

Taught by your **amazing TAs** and will

• recap lecture
• show you **how to apply** the ideas you learned in lecture
• can occasionally cover new material

Sections are as “mandatory” as lectures:

• we will not track attendance, but
• sections (practice, practice, practice) are the **best** way to learn the material in CS 161
• also, a good place to find community
Homework

• Weekly assignments, posted Wednesday by 11:59pm, due the next Wednesday 11:59pm.
• First HW will be posted this Wednesday
• There are 8 total (no HW the week before midterm)
  • HW1, HW2, HW3: solo submissions
  • HW4, HW5, HW6, HW7, HW8: solo/pair submissions
How to get the most out of homework

• HW has two parts: exercises and problems.
• Do the exercises on your own.
• Try the problems on your own before discussing it with classmates.

• If you get help from a CA during office hours:
  • Try the problem first.
  • Ask: “I was trying this approach and I got stuck here.”
  • After you’ve figured it out, write up your solution from scratch, without the notes you took during office hours.
Exams

• There will be a **midterm** and a **final**
  • **Midterm:** Thu Feb 15, 6:00pm–9:00pm
    • Covers lectures 1-7
  • **Final:** Mon Mar 18, 3:30pm-6:30pm
    • Covers everything, but more focus on lectures 8 onwards

• Ensure you can show up for both.

• If you have a conflict with the midterm time, email cs161-staff-win2324@cs.stanford.edu ASAP!!!!

• We cannot accommodate final exam conflicts.
Grading

• Homework 50%
  • Lowest homework score will be dropped. Each of the 7 remaining homeworks equally weighted.
• Midterm 20%
• Final 30%

• Letter grade thresholds not fixed, and will be based on what instructors determine at the end showed excellent, very good, good, etc. performance
Talk to us!

• Ed discussion forum:
  • Link on top of the course website
  • Course announcements will be posted there
  • Discuss material with TAs and your classmates

• Office hours:
  • See course website for schedule
  • Some online, some in-person
  • They start later this week
Talk to each other!

- Answer your peers’ questions on Ed!
Course elements and resources

• Course website:
  • cs161.stanford.edu

• Lectures
  • Pre-Lecture Exercises
  • Lecture Notes
  • IPython Notebooks
  • Concept Check Questions

• References
• Sections
• Homework
• Exams
• Office Hours and Ed
A note on course policies

• Course policies are listed on the website.
• Read them and adhere to them.
• That’s all I’m going to say about course policies (except for a couple of slides on collaboration and the honor code)
Collaboration

• We encourage collaboration on homework (but strongly recommend you do exercises on your own)

• Valid and invalid modes of collaboration are detailed on the course website.
  
  • Briefly, you can exchange ideas with classmates but must write up solutions on your own (except for pair submissions from HW4 onwards).
  
  • For pair submissions (HW4 onwards), each person must understand and contribute to the solution for each individual problem.

• You must cite all collaborators, as well as all sources used (outside of course materials).
Honor code

• Updated two years ago: “In all cases, it is not permissible for students to enter exam questions into any software, apps, or websites. Accessing resources that directly explain how to answer questions from the actual assignment or exam is a violation of the Honor Code.”

• Course policy for homework: “In all cases, it is not permissible for students to enter homework questions into any software, apps, or websites. Accessing resources that directly explain how to answer questions from the actual assignment or exam is a violation of course policy.”
Bug bounty!

• We hope all PSETs and slides will be bug-free.

• However, we sometimes make typos.

• **If you find a typo** (that affects understanding*) on slides, IPython notebooks, Section material or PSETs:
  • Let us know! (Post on Ed or tell a CA).
  • The first person to catch a serious bug might get a good citizenship bonus point.

*So, typos *ike thees onse don’t count, although please point those out too. Typos like $2 + 2 = 5$ do count, as does pointing out that we omitted some crucial information.
For SCPD Students (and all students)

• Some office hours held online
• One of the recitation sections will be recorded.
• See the website for more details! (coming soon...)
OAE forms

• Please send OAE forms to

  cs161-staff-win2324@cs.stanford.edu
Feedback!

• We will have high-resolution feedback throughout the course (subset of you randomly asked each week, starting week 2).
• Please help us improve the course!
How are you approaching CS 161?
Everyone can succeed in this class!

1. Work hard
2. Work smart
3. Ask for help
The big questions

• Who are we?
  • Professor, TA’s, students?

• Why are we here?
  • Why learn about algorithms?

• What is going on?
  • What is this course about?
  • Logistics?
  • Embedded Ethics?

• Can we multiply integers?
  • And can we do it quickly?
Welcome to Embedded EthiCS

Dan Webber, PhD (he/him)
webberdf@stanford.edu
Hi! I’m Dan Webber

• BA in Computer Science, Amherst College
  • Afterwards: a couple years as a software developer in fintech and e-commerce
• PhD in Philosophy, University of Pittsburgh
  • Moral theory and social/political philosophy
• Now: Postdoc, EIS and HAI at Stanford
  • Embedding ethics into CS courses like this one!
What is Embedded EthiCS?

Training the next generation of computer scientists to “consider ethical issues from the outset rather than building technology and letting problems surface downstream” by integrating skills and habits of ethical analysis throughout the Stanford Computer Science curriculum.

Elan the Ethical Emu
<table>
<thead>
<tr>
<th>Institution</th>
<th>Course Numbers</th>
</tr>
</thead>
<tbody>
<tr>
<td>Harvard (2017)</td>
<td>CS106A</td>
</tr>
<tr>
<td>Georgetown (2017)</td>
<td>CS106B</td>
</tr>
<tr>
<td>Brown (2019)</td>
<td>CS107</td>
</tr>
<tr>
<td>Northeastern (2019)</td>
<td>CS109</td>
</tr>
<tr>
<td>MIT (2020)</td>
<td>CS111</td>
</tr>
<tr>
<td>... and many other places</td>
<td></td>
</tr>
<tr>
<td></td>
<td>CS147</td>
</tr>
<tr>
<td></td>
<td>CS161</td>
</tr>
<tr>
<td></td>
<td>CS177</td>
</tr>
<tr>
<td></td>
<td>CS221</td>
</tr>
<tr>
<td></td>
<td>CS224n</td>
</tr>
<tr>
<td></td>
<td>CS234</td>
</tr>
<tr>
<td></td>
<td>CS247B</td>
</tr>
<tr>
<td></td>
<td>soon: CS107E and more!</td>
</tr>
</tbody>
</table>
What do we teach?

- **Issue spotting** and ethical sensitivity.
- Recognizing values in **design choices**.
- Developing **language** to talk about moral choices.
- **Professional responsibilities** of computer scientists & software engineers.
- Important topics in **technology ethics**: bias & fairness, inequity, privacy, surveillance, data control & consent, trust, disinformation, participatory design, concentration of power.
Our guiding questions
• Does it work?
• Is it fast?
• Can I do it better?
• Can I do it more ethically?
• (Should I do it all?)
Catch you soon!

Dan Webber
webberdf@stanford.edu

Want to talk more about ethics? Email to set up a meeting!
The big questions

- Who are we?
  - Professor, TA’s, students?
- Why are we here?
  - Why learn about algorithms?
- What is going on?
  - What is this course about?
  - Logistics?
  - Embedded Ethics?
- Can we multiply integers?
  - And can we do it quickly?
Course goals

- Think analytically about algorithms
- Flesh out an “algorithmic toolkit”
- Learn to communicate clearly about algorithms

Today’s goals

- Karatsuba Integer Multiplication
- Algorithmic Technique:
  - Divide and conquer
- Algorithmic Analysis tool:
  - Intro to asymptotic analysis
Let’s start at the beginning
Etymology of “Algorithm”

- Al-Khwarizmi was a 9th-century scholar, born in present-day Uzbekistan, who studied and worked in Baghdad during the Abbassid Caliphate.
- Among many other contributions in mathematics, astronomy, and geography, he wrote a book about how to multiply with Arabic numerals.
- His ideas came to Europe in the 12th century.

* Dixit algorizmi (so says Al-Khwarizmi)

- Originally, “Algorisme” [old French] referred to just the Arabic number system, but eventually it came to mean “Algorithm” as we know today.
This was kind of a big deal

\[ \text{XLIV} \times \text{XCVII} = ? \]

\[ \begin{array}{c}
44 \\
\times 97 \\
\end{array} \]
Integer Multiplication

44
\times 97

___
Integer Multiplication

1234567895931413

\times 4563823520395533

\hline


How fast is the grade-school multiplication algorithm?

(How many one-digit operations?)

About $n^2$ one-digit operations

At most $n^2$ multiplications, and then at most $n^2$ additions (for carries) and then I have to add $n$ different 2n-digit numbers...
Big-Oh Notation

• We say that Grade-School Multiplication “runs in time $O(n^2)$”

• Formal definition coming Wednesday!
• Informally, big-Oh notation tells us how the running time scales with the size of the input.
Implemented in Python, on a laptop

The runtime “scales like” $n^2$

Looks like it’s roughly

$$T_{\text{laptop}}(n) = 0.0063 n^2 - 0.5 n + 12.7 \text{ ms}...$$
Implemented by hand

The runtime still “scales like” $n^2$

$$T_{\text{laptop}}(n) \approx 0.0063 n^2 - 0.5 n + 12.7 \text{ ms}$$
Why is big-Oh notation meaningful?

\[ \approx \frac{n^{1.6}}{10} + 100 \]

\[ \approx .0063n^2 \]
Let $n$ get bigger...

\[ \approx \frac{n^{1.6}}{10} + 100 \]

\[ \approx .0063n^2 \]
Take-away

• An algorithm that runs in time $O(n^{1.6})$ is “better” than an algorithm that runs in time $O(n^2)$.

• So the question is...
Can we do better?

Can we multiply n-digit integers faster than $O(n^2)$?
Let’s dig into our algorithmic toolkit...
Divide and conquer

Break problem up into smaller (easier) sub-problems

- Big problem
  - Smaller problem
    - Yet smaller problem
    - Yet smaller problem
  - Smaller problem
    - Yet smaller problem
    - Yet smaller problem
    - Yet smaller problem

Divide and conquer for multiplication

Break up an integer:

\[ 1234 = 12 \times 100 + 34 \]

\[ 1234 \times 5678 \]
\[ = (12 \times 100 + 34) (56 \times 100 + 78) \]
\[ = (12 \times 56)10000 + (34 \times 56 + 12 \times 78)100 + (34 \times 78) \]

One 4-digit multiply

Four 2-digit multiplies
More generally

Break up an n-digit integer:

\[
[x_1x_2\cdots x_n] = [x_1x_2\cdots x_{n/2}] \times 10^{n/2} + [x_{n/2+1}x_{n/2+2}\cdots x_n]
\]

\[
x \times y = (a \times 10^{n/2} + b)(c \times 10^{n/2} + d)
\]

\[
= (a \times c)10^n + (a \times d + c \times b)10^{n/2} + (b \times d)
\]

One n-digit multiply \[\rightarrow\] Four (n/2)-digit multiplies
Divide and conquer algorithm

x, y are n-digit numbers

(Multiply) (x, y):

• If n=1:
  • Return xy

• Write \( x = a \ 10^{\frac{n}{2}} + b \)

• Write \( y = c \ 10^{\frac{n}{2}} + d \)

• Recursively compute \( ac, ad, bc, bd \):  
  • \( ac = \text{Multiply}(a, c) \), etc.. 

• Add them up to get \( xy \):
  • \( xy = ac \ 10^n + (ad + bc) \ 10^{n/2} + bd \)

Base case: I’ve memorized my 1-digit multiplication tables...

a, b, c, d are n/2-digit numbers

Make this pseudocode more detailed! How should we handle odd n? How should we implement “multiplication by 10^n”?

See the Lecture 1 Python notebook for actual code!
Think-Pair-Share

• We saw that this 4-digit multiplication problem broke up into four 2-digit multiplication problems

1234 × 5678

• If you recurse on those 2-digit multiplication problems, how many 1-digit multiplications do you end up with in total?
Recursion Tree

4 digits

2 digits

2 digits

2 digits

2 digits

16 one-digit multiplies!
What is the running time?

• Better or worse than the grade school algorithm?

• How do we answer this question?
  1. Try it.
  2. Try to understand it analytically.
1. Try it.

Check out the Lecture 1 IPython Notebook

Conjectures about running time?

Doesn’t look too good but hard to tell…

Maybe one implementation is slicker than the other?

Maybe if we were to run it to n=10000, things would look different.

Something funny is happening at powers of 2…
2. Try to understand the running time analytically

• Proof by meta-reasoning:

  It must be faster than the grade school algorithm, because we are learning it in an algorithms class.

Not sound logic!
2. Try to understand the running time analytically

Think-Pair-Share:

• We saw that multiplying 4-digit numbers resulted in 16 one-digit multiplications.

• How about multiplying 8-digit numbers?

• What do you think about n-digit numbers?
Recursion Tree

8 digits

4 digits

4 digits

4 digits

4 digits

64 one-digit multiplies!
2. Try to understand the running time analytically

Claim:

The running time of this algorithm is AT LEAST $n^2$ operations.
There are $n^2$ 1-digit problems

1 problem of size $n$

4 problems of size $n/2$

...$4^t$ problems of size $n/2^t$

Note: this is just a cartoon – I’m not going to draw all $4^t$ circles!

- If you cut $n$ in half $\log_2(n)$ times, you get down to 1.
- So at level $t = \log_2(n)$ we get...

\[ 4^{\log_2 n} = n^{\log_2 4} = n^2 \]

$n^2$ problems of size 1

...
That’s a bit disappointing
All that work and still (at least) $O(n^2)$...

But wait!!
Divide and conquer can actually make progress

- Karatsuba figured out how to do this better!

\[
xy = (a \cdot 10^{n/2} + b)(c \cdot 10^{n/2} + d)
\]

\[
= ac \cdot 10^n + (ad + bc)10^{n/2} + bd
\]

Need these three things

- If only we could recurse on three things instead of four...
Karatsuba integer multiplication

- Recursively compute these THREE things:
  - $ac$
  - $bd$
  - $(a+b)(c+d)$

  $(a+b)(c+d) = ac + bd + bc + ad$

- Assemble the product:

$$xy = (a \cdot 10^{n/2} + b)(c \cdot 10^{n/2} + d)$$

$$= ac \cdot 10^n + (ad + bc)10^{n/2} + bd$$
How would this work?

x, y are n-digit numbers

**Multiply**(x, y):

- If n=1:
  - Return xy
- Write \( x = a \cdot 10^{\frac{n}{2}} + b \) and \( y = c \cdot 10^{\frac{n}{2}} + d \)
- ac = **Multiply**(a, c)
- bd = **Multiply**(b, d)
- z = **Multiply**(a+b, c+d)
- \( xy = ac \cdot 10^n + (z - ac - bd) \cdot 10^{n/2} + bd \)
- Return xy

(Still not super precise, see IPython notebook for detailed code. Also, still assume n is a power of 2.)

a, b, c, d are n/2-digit numbers
What’s the running time?

1 problem of size n

3 problems of size n/2

3\textsuperscript{t} problems of size n/2\textsuperscript{t}

- If you cut n in half \( \log_2(n) \) times, you get down to 1.
- So at level \( t = \log_2(n) \) we get...

\[ 3^{\log_2 n} = n^{\log_2 3} \approx n^{1.6} \]

problems of size 1.

Note: this is just a cartoon – I’m not going to draw all \( 3^t \) circles!

\( n^{1.6} \) problems of size 1

We aren’t accounting for the work at the higher levels!
But we’ll see later that this turns out to be okay.
This is much better!
We can even see it in real life!

**Multiplying n-digit integers**

- **Grade School Multiplication**
- **Divide and Conquer II (Karatsuba)**

---

The graph compares the time (in milliseconds) taken to multiply n-digit integers using two methods: Grade School Multiplication and Divide and Conquer II (Karatsuba). As the number of digits (n) increases, the time taken for Grade School Multiplication (red line) grows exponentially, while Divide and Conquer II (orange line) grows at a slower rate, demonstrating the efficiency of the latter method.
Can we do better?

- **Toom-Cook** (1963): instead of breaking into three $n/2$-sized problems, break into five $n/3$-sized problems.
  - Runs in time $O(n^{1.465})$

  Try to figure out how to break up an $n$-sized problem into five $n/3$-sized problems! *(Hint: start with nine $n/3$-sized problems).*

- **Schönhage–Strassen** (1971): 
  - Runs in time $O(n \log(n) \log \log(n))$

- **Furer** (2007) 
  - Runs in time $n \log(n) \cdot 2^{O(\log^*(n))}$

- **Harvey and van der Hoeven** (2019) 
  - Runs in time $O(n \log(n))$

  [This is just for fun, you don’t need to know these algorithms!]
Course goals

• Think analytically about algorithms
• Flesh out an “algorithmic toolkit”
• Learn to communicate clearly about algorithms

Today’s goals

• Karatsuba Integer Multiplication
• Algorithmic Technique:
  • Divide and conquer
• Algorithmic Analysis tool:
  • Intro to asymptotic analysis
How was the pace today?
The big questions

• Who are we?
  • Professor, TA’s, students?

• Why are we here?
  • Why learn about algorithms?

• What is going on?
  • What is this course about?
  • Logistics?

• Can we multiply integers?
  • And can we do it quickly?

• Wrap-up
Wrap up

• cs161.stanford.edu

• Algorithms are fundamental, useful and fun!

• In this course, we will develop both algorithmic intuition and algorithmic technical chops

• Karatsuba Integer Multiplication:
  • You can do better than grade school multiplication!
  • Example of divide-and-conquer in action
  • Informal demonstration of asymptotic analysis
Next time

• Sorting!
• Asymptotics and (formal) Big-Oh notation
• Divide and Conquer some more

BEFORE Next time

• Pre-lecture exercise!  On the course website!