

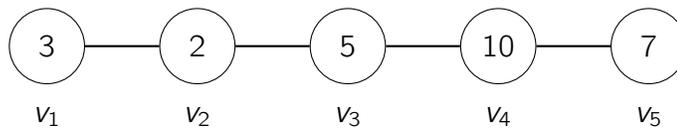
1 LCS Warm-up

- Given sequences $X = ABCB$ and $Y = BDCAB$, fill in the LCS dynamic programming table $C[i, j]$ below, where $C[i, j]$ is the length of the LCS of $X[1 : i]$ and $Y[1 : j]$. Then use the completed table to recover an actual longest common subsequence.

	\emptyset	B	D	C	A	B
\emptyset	0	0	0	0	0	0
A	0					
B	0					
C	0					
B	0					

2 DP Exercises

- Consider a path graph with 5 vertices, where each vertex v_i has weight w_i as shown. An **independent set** is a subset of vertices such that no two are adjacent. Find a maximum weight independent set.



- Define $A(i)$ to be the weight of the maximum weight independent set considering only vertices v_1, \dots, v_i . Write a recurrence for $A(i)$.
 - Compute $A(i)$ for $i = 0, 1, \dots, 5$. Which vertices are in the optimal set?
 - What is the runtime of the resulting algorithm?
- In how many ways can you tile a $2 \times n$ board using 1×2 dominoes (placed either horizontally or vertically)?
 - Let $T(n)$ be the number of tilings of a $2 \times n$ board. Write a recurrence for $T(n)$ and justify it.

[Hint: Consider how the rightmost column can be filled.]
 - Compute $T(n)$ for $n = 1, 2, \dots, 7$.
 - Does this recurrence look familiar?

3 Activity Selection: Greedy Strategies

The activity selection problem asks: given n activities with start and finish times, find a maximum-size set of non-conflicting activities. In lecture, we saw that always selecting the activity with the earliest **finish time** (that doesn't conflict) gives an optimal solution.

For each of the following alternative greedy strategies, give a counterexample showing it does not always produce an optimal solution.

1. Always select the activity with the earliest **start time** that doesn't conflict with already-selected activities.
2. Always select the activity with the **shortest duration** that doesn't conflict with already-selected activities.

4 Minimum Refueling Stops

You need to drive from city A to city B , a total distance of L miles. Your car can travel at most D miles on a full tank. Gas stations are located at distances $g_1 < g_2 < \dots < g_n$ from city A . You start with a full tank. You are guaranteed that the distance between any two consecutive stops is at most D (where city A and city B are also considered stops), so it is always possible to reach city B .

Design an algorithm that finds the minimum number of refueling stops needed to reach city B .

- (a) Describe a greedy strategy for this problem.
- (b) Apply your strategy with $L = 20$, $D = 7$, and stations at positions 3, 5, 8, 12, 14, 17. List the positions where you stop to refuel.
- (c) Prove your greedy strategy is optimal.

[Hint: Show by induction that after k stops, the greedy car is always at least as far along as in any other solution.]